

(19) World Intellectual Property Organization
International Bureau



(43) International Publication Date
8 March 2001 (08.03.2001)

PCT

(10) International Publication Number
WO 01/17252 A1

(51) International Patent Classification⁷: H04N 7/167

(21) International Application Number: PCT/US00/22834

(22) International Filing Date: 18 August 2000 (18.08.2000)

(25) Filing Language: English

(26) Publication Language: English

(30) Priority Data:
09/385,590 29 August 1999 (29.08.1999) US

(71) Applicant (for all designated States except US): INTEL CORPORATION [US/US]; 2200 Mission College Boulevard, Santa Clara, CA 95052 (US).

(72) Inventors; and

(75) Inventors/Applicants (for US only): GRAUNKE, Gary, L. [US/US]; 362 NE Hillwood Drive, Hillsboro, OR 97124 (US). LEE, David, A. [US/US]; 740 SW Willow Creek Drive, Beaverton, OR 97006 (US). FABER, Robert, W. [US/US]; 942 NE Third Avenue, Hillsboro, OR 97124 (US). TRAW, Brendan, S. [US/US]; 10859 NW Supreme Court, Portland, OR 97229 (US).

(74) Agents: MALLIE, Michael, J. et al.; Blakely, Sokoloff, Taylor & Zafman LLP, 7th floor, 12400 Wilshire Boulevard, Los Angeles, CA 90025 (US).

(81) Designated States (national): AE, AG, AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, BZ, CA, CH, CN, CR, CU, CZ, DE, DK, DM, DZ, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MA, MD, MG, MK, MN, MW, MX, MZ, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, TZ, UA, UG, US, UZ, VN, YU, ZA, ZW.

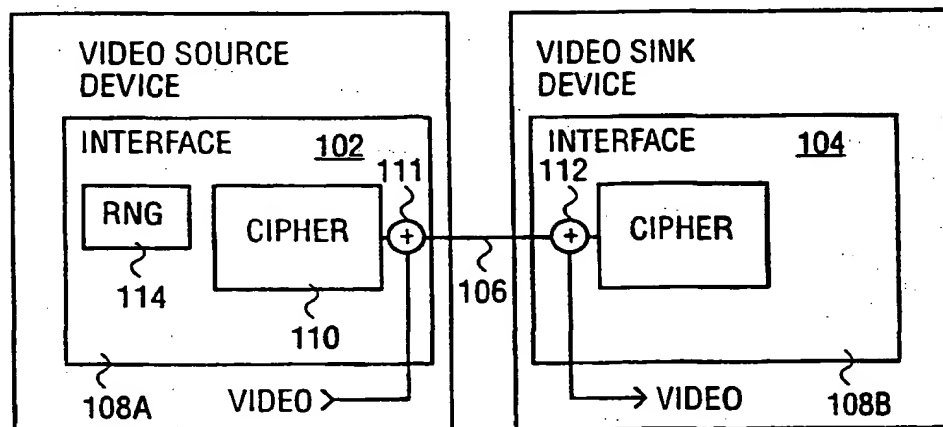
(84) Designated States (regional): ARIPO patent (GH, GM, KE, LS, MW, MZ, SD, SL, SZ, TZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).

Published:

— With international search report.

For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.

(54) Title: DIGITAL VIDEO CONTENT TRANSMISSION CIPHERING AND DECIPHERING METHOD AND APPARATUS



(57) Abstract: A video source device provides a basis value to a symmetric ciphering/deciphering process to a video sink device, to which the video source device is to provide a video content. The video source device ciphers the video content for transmission to the video sink device, including generation of a first cipher key through functional transformation of the basis value. The video source device further verifies periodically that the transmitted video content is indeed being symmetrically deciphered by the video sink device. The video sink device practices symmetric deciphering, including generation of a first decipher key through functional transformation of the basis value. The video sink device also provides verification values to the video source device to facilitate confirmation of symmetric deciphering. In one embodiment, the video source and sink devices further authenticate each other, including generation of an authentication key, which is used in the generation of the first cipher/decipher key.

THIS PAGE BLANK (USPTO)

Digital Video Content Transmission Ciphering And Deciphering
Method And Apparatus

BACKGROUND OF THE INVENTION

1. **Field of the Invention**

The present invention relates to the field of content protection. More specifically, the present invention addresses the provision of protection to digital video content to facilitate their secure transmission from a video source device to a video sink device.

2. **Background Information**

In general, entertainment, education, art, and so forth (hereinafter collectively referred to as "content") packaged in digital form offer higher audio and video quality than their analog counterparts. However, content producers, especially those in the entertainment industry, are still reluctant in totally embracing the digital form. The primary reason being digital contents are particularly vulnerable to pirating. As unlike the analog form, where some amount of quality degradation generally occurs with each copying, a pirated copy of digital content is virtually as good as the "gold master". As a result, much effort have been spent by the industry in developing and adopting techniques to provide protection to the distribution and rendering of digital content.

Historically, the communication interface between a video source device (such as a personal computer) and a video sink device (such as a monitor) is an analog interface. Thus, very little focus has been given to providing protection for the transmission between the source and sink devices. With advances in integrated circuit and other related technologies, a new type of digital interface between video source and sink devices is emerging. The availability of this type of new digital interface presents yet another new challenge to protecting digital video content. While in general, there is a large body of cipher technology known, the operating characteristics such as the volume of the data, its

streaming nature, the bit rate and so forth, as well as the location of intelligence, typically in the source device and not the sink device, present a unique set of challenges, requiring a new and novel solution.

SUMMARY OF THE INVENTION

A video source device provides a basis value to a symmetric ciphering/deciphering process to a video sink device, to which the video source device is to provide a video content. The video source device ciphers the video content for transmission to the video sink device, including generation of a first cipher key through functional transformation of the basis value. The video source device further verifies periodically that the transmitted video content is indeed being symmetrically deciphered by the video sink device.

BRIEF DESCRIPTION OF DRAWINGS

The present invention will be described by way of exemplary embodiments, but not limitations, illustrated in the accompanying drawings in which like references denote similar elements, and in which:

Figure 1 illustrates an overview of the present invention in accordance with one embodiment;

Figure 2 illustrates a symmetric ciphering/deciphering process based method for providing video content from a source device to a sink device, in accordance with one embodiment;

Figures 3a-3b illustrate the symmetric ciphering/deciphering process of **Fig. 2**, in accordance with one embodiment;

Figure 4 illustrates video source and sink devices of **Fig. 1** in further detail, in accordance with one embodiment;

Figure 5 illustrates the combined block/stream cipher of **Fig. 4** in further detail, in accordance with one embodiment;

Figure 6 illustrates the block key section of **Fig. 5** in further detail, in accordance with one embodiment;

Figure 7 illustrates the block data section of **Fig. 5** in further detail, in accordance with one embodiment; and

Figures 8a-8c illustrate the stream data section of **Fig. 5** in further detail, in accordance with one embodiment.

DETAILED DESCRIPTION OF THE INVENTION

In the following description, various aspects of the present invention will be described, and various details will be set forth in order to provide a thorough understanding of the present invention. However, it will be apparent to those skilled in the art that the present invention may be practiced with only some or all aspects of the present invention, and the present invention may be practiced without the specific details. In other instances, well known features are omitted or simplified in order not to obscure the present invention.

Various operations will be described as multiple discrete steps performed in turn in a manner that is most helpful in understanding the present invention. However, the order of description should not be construed as to imply that these operations are necessarily performed in the order they are presented, or even order dependent. Lastly, repeated usage of the phrase "in one embodiment" does not necessarily refer to the same embodiment, although it may.

Referring now to **Figure 1**, wherein a block diagram illustrating an overview of the present invention, in accordance with one embodiment is shown. As illustrated, video source device **102** and video sink device **104** are coupled to each other by digital video link **106**. Video source device **102** provides video content to video sink device **104** through digital video link **106**. In accordance with the present invention, video source device **102** and video sink device **104** are equipped to be able to jointly practice a symmetric ciphering/deciphering process. As a result, video content may be provided in a more robust ciphered digital form from video source device **102** to video sink device **104** through video link **106**, making it more difficult to pirate video content during transmission.

Except for the teachings of the present invention incorporated, to be described more fully below, video source device **102** and video sink device **104** are both intended to represent a broad range of such devices known in the art. Examples of video source devices include but not limited to computers of all sizes (from palm size device to desktop device, and beyond), set-up boxes, or DVD players, whereas examples of video sink devices include but not limited to CRT monitors, flat panel displays or television sets. Digital video link **106** may be implemented in any one of a number of mechanical and electrical forms, as long as they are consistent with the operating requirement (i.e. speed, bit rate and so forth), and a mechanism (which may be in hardware or through protocol) is provided to allow control information to be exchanged between video source and sink devices **102** and **104** (hereinafter, simply source and sink devices respectively).

Figure 2 illustrates an overview of the symmetric ciphering/deciphering process based method for providing video content from a source device to a sink device, in accordance with one embodiment. In this embodiment, source and sink devices **102** and **104** are assumed to have each been provided with an array of private keys and a complementary identifier by a certification authority. As illustrated, upon power on or reset, source device **102** first provides a basis value to the symmetric ciphering/deciphering process to sink device **104** (block **202**). For the illustrated embodiment, the basis value is a random number (A_n). A_n may be generated in any one of a number of techniques known in the art. Additionally, source device **102** also provides its identifier (A_k) to sink device **104** (block **202**). In response, sink device **104** replies with its identifier (B_k) (block **203**). Upon exchanging the above information, source and sink devices **102** and **104** independently generate their respective copies of an authentication key (K_m) using A_k and B_k (block **204** and **205**). For the illustrated embodiment, source device **102** generates its copy of K_m by summing private keys of its provided array indexed by B_k , while sink device **104** generates its copy of K_m by summing private keys of its provided array indexed by A_k . At this time, if both source and

sink devices **102** and **104** are authorized devices, they both possess and share a common secret authentication key K_m .

In one embodiment, each of source and sink devices **102** and **104** is pre-provided with an array of 40 56-bit private keys by the certification authority. A_n is a 64-bit random number, and K_m is 56-bit long. For more information on the above described authentication process, see co-pending U.S. Patent Application, serial number 09/275,722, filed on March 24, 1999, entitled Method and Apparatus for the Generation of Cryptographic Keys, having common inventorship as well as assignee with the present application.

Having authenticated sink device **104**, source device **102** ciphers video content into a ciphered form before transmitting the video content to sink device **104**. Source device **102** ciphers the video content employing a symmetric ciphering/deciphering process, and using the random number (A_n) as well as the independently generated authentication key (K_m) (block **206**). Upon receipt of the video content in ciphered form, sink device **104** deciphers the ciphered video content employing the same symmetric ciphering/deciphering processing, and using the provided A_n as well as its independently generated copy of K_m (block **207**).

In accordance with the present invention, as an integral part of ciphering video content, source device **102** derives a set of verification reference values in a predetermined manner (block **208**). Likewise, as an integral part of symmetrically deciphering video content sink device **104** also derives a set of verification values in a predetermined manner, and transmits these derived verification values to source device **102** (block **209**). Upon receiving each of these verification values, source device **102** compares the received verification value to the corresponding one of the verification reference value to determine and confirm that indeed the ciphered video content is being properly deciphered by sink device **104** (block **210**).

For the illustrated embodiment, both source and sink devices **102** and **104** generate the verification reference and verification values continuously, but the verification values are provided from sink device **104** to source device **102** periodically at predetermined intervals.

In one embodiment, the verification reference and verification values are all 64-bit in length, and sink device **104** provides source device **102** with verification values at initialization and every 64th frames thereafter.

Figures 3a-3b illustrate the symmetric ciphering/deciphering process in further detail, in accordance with one embodiment. In this embodiment, the video content is assumed to be a multi-frame video content with each frame having multiple lines of video content. In between two lines of a frame is an interval to allow a sink device to horizontally "retrace" itself, commonly known as the horizontal retrace interval or horizontal blanking interval (HBI). Likewise, in between two frames is an interval to allow a sink device to vertically "retrace" itself, commonly known as the vertical retrace interval or vertical blanking interval (VBI).

Source device **102** first generates a session key (K_s) for the transmission session (block **302**). For the illustrated embodiment, K_s is generated by block ciphering the above mentioned random number A_n using the authentication key K_m as the block cipher key and applying C_1 clocks. The duration of a transmission session is application dependent. Typically, it corresponds to a natural demarcation of the video content, e.g. the transmission of a single movie may constitute a transmission session, or the transmission of an episode of a sitcom may constitute a transmission session instead.

Upon generating the session key K_s , source device **102** generates an initial version of a second random number (M_0) (block **304**). For the illustrated embodiment, source device **102** first generates a pseudo random bit sequence (at p -bit per clock) using a stream cipher with the above described random number A_n and the session key K_s (in two roles, as another input random number and as the stream cipher key), applying C_2 clocks. Source device **102** derives M_0 from the pseudo random bit sequence, as the bit sequence is generated.

Next, source device **102** generates a frame key (K_i) for the next frame (block **306**). For the illustrated embodiment, K_i is generated by block ciphering an immediately preceding version of the second random number M_{i-1} using the

session key K_s as the block cipher key, and applying C_3 clocks. That is, for the first frame, frame-1, frame key K_1 is generated by block ciphering the above described initial version of the second random number M_0 , using K_s , and applying C_3 clocks. Additionally, this operation is subsequently repeated at each vertical blanking interval for the then next frame, frame-2, frame-3, and so forth.

Upon generating the frame key K_i , source device 102 generates the current version of the second random number (M_i) (block 302). For the illustrated embodiment, source device 102 first generates a pseudo random bit sequence (at p-bit per clock) using a stream cipher with the previous version of the second random number M_{i-1} and the frame key K_i (in two roles, as another input random number and as the stream cipher key), applying C_4 clocks. Source device 102 derives M_i from the pseudo random bit sequence, as the bit sequence is generated.

Upon generating the current version of the second random number M_i , source device 102 again generates a pseudo random bit sequence (at p-bit per clock) to cipher the frame (block 308). For the illustrated embodiment, source device 102 generates the pseudo random bit sequence using a stream cipher with an immediately preceding version of the second random number M_{i-1} and frame key K_i (in two roles, as another input random number and the stream cipher key), applying C_5 clock cycles. The video content is ciphered by performing an exclusive-OR (XOR) operation on the video stream and the pseudo random bit sequence. The pseudo random bit sequence is generated preferably at a rate sufficient to cipher a pixel of RGB signal per clock. Therefore, C_5 is equal to the number of bits per pixel multiply by the number of pixels per line, as well as the number of lines per frame.

For the illustrated embodiment, a stream cipher that successively transforms M_i and K_i in the course of generating the pseudo random bit sequence is employed. Additionally, the robustness of the ciphered video content is further strengthened by increasing the unpredictability of the pseudo random bit sequence through successive modification of the current states of K_i at the horizontal blanking intervals of the frame (block 310).

Over in sink device **104**, in like manner, it first generates a session key (Ks) for the transmission session (block **312**). Upon generating the session key Ks, sink device **104** generates an initial version of the second random number (M0) (block **314**). Next, sink device **104** generates the frame key (Ki) and second random number (Mi) for the next frame (block **316**). This operation is likewise subsequently repeated at each vertical blanking interval for the then next frame. In the meantime, after generation of each frame key Ki and Mi, sink device **104** generates a corresponding pseudo random bit sequence to decipher the frame (block **318**). The ciphered video content is deciphered by performing an exclusive-OR (XOR) operation on the video stream and the corresponding pseudo random bit sequence. Sink device **104** also employs a stream cipher that successively transforms Mi and Ki in the course of generating the pseudo random bit sequence. Furthermore, Ki is successively modified at the horizontal blanking intervals of the frame (block **320**). Ki, the pseudo random bit sequence, and Mi are symmetrically generated as earlier described for source device **102**.

In one embodiment, Ks and each Ki are both 84-bit in length. C1 and C3 are both 48 clocks in length. Each pixel is 24-bit, and the pseudo random bit sequence is generated at 24-bit per clock. Each Mi is 64-bit in length, C3 and C4 are 56 clocks in length. Each 64-bit Mi is formed by concatenating the "lower" 16-bit stream cipher output of each of the last four clocks.

Accordingly, video content may be advantageously transmitted in ciphered form with increased robustness from source device **102** to sink device **104** through link **106** with reduced pirating risk.

Figure 4 illustrates video source and sink devices of **Fig. 1** in further detail, in accordance with one embodiment. As shown, video source and sink devices **102** and **104** include interfaces **108a** and **108b** disposed at the respective end of link **106**. Each of interfaces **108a** and **108b** is advantageously provided with cipher **110** of the present invention and XOR **112** to practice the video content protection method of the present invention as described above. Additionally, for ease of explanation, interface **108a** is also shown as having been provided with a separate random number generator **114**. Except for

interfaces **108a** and **108b**, as stated earlier, video source and sink devices **102** and **104** are otherwise intended to represent a broad category of these devices known in the art.

Random number generator **114** is used to generate the earlier described random number A_n . Random number generator **114** may be implemented in hardware or software, in any one of a number of techniques known in the art. In alternate embodiments, as those skilled in the art will appreciate from the description to follow, cipher **110** may also be used to generate A_n , without the employment of a separate random number generator.

Cipher **110** is a novel combined block/stream cipher capable of operating in either a block mode of operation or a stream mode of operation. To practice the video content protection method of the present invention, cipher **110** is used in block mode to generate the above described session key K_s and frame keys K_i , and in stream mode to generate the pseudo random bit sequences for the various frames (and indirectly M_i , as they are derived from the respective bit sequences).

In source device **102**, XOR **112** is used to cipher video content, combining it with the pseudo random bit sequences generated by cipher **110** on interface **108a**. Over in sink device **104**, XOR **112** is used to decipher ciphered video content, combining it with the pseudo random bit sequences generated by cipher **110** on interface **108b**.

Figure 5 illustrates the combined block/stream cipher of Fig. 4 in further detail, in accordance with one embodiment. As illustrated, combined block/stream cipher **110** includes block key section **502**, data section **504**, stream key section **506**, and mapping section **508**, coupled to one another. Block key section **502** and data section **504** are employed in both the block mode as well as the stream mode of operation, whereas stream key section **506** and mapping section **508** are employed only in the stream mode of operation.

Briefly, in block mode, block key section **502** is provided with a block cipher key, such as the earlier described authentication key K_m or the session key K_s ; whereas data section **504** is provided with the plain text, such as the

earlier described random number A_n or the derived random number M_{i-1} .

"Rekeying enable" signal is set to a "disabled" state, operatively de-coupling block key section 502 from stream key section 506. During each clock cycle, the block cipher key as well as the plain text are transformed. The block cipher key is independently transformed, whereas transformation of the plain text is dependent on the transformation being performed on the block cipher key. After a desired number of clock cycles, the provided plain text is transformed into ciphered text. For the earlier described video content protection method, when block key section 502 is provided with K_m and data section 504 is provided with the A_n , ciphered A_n is read out and used as the session key K_s . When block key section 502 is provided with K_s and data section 504 is provided with the M_{i-1} , ciphered M_{i-1} is read out and used as the frame key K_i .

To decipher the ciphered plain text, block key section 502 and data section 504 are used in like manner as described above to generate the intermediate "keys", which are stored away (in storage locations not shown). The stored intermediate "keys" are then applied to the ciphered text in reversed order, resulting in the deciphering of the ciphered text back into the original plain text. Another approach to deciphering the ciphered text will be described after block key section 502 and data section 504 have been further described in accordance with one embodiment each, referencing Figs. 6-7.

In stream mode, stream key section 506 is provided with a stream cipher key, such as the earlier described session key K_s or frame key K_i . Block key section 502 and data section 504 are provided with random numbers, such as the earlier described session/frame keys K_s/K_i and the derived random numbers M_{i-1} . "Rekeying enable" signal is set to an "enabled" state, operatively coupling block key section 502 to stream key section 506. Periodically, at predetermined intervals, such as the earlier described horizontal blanking intervals, stream key section 506 is used to generate one or more data bits to dynamically modify the then current state of the random number stored in block data section 502. During each clock cycle, in between the predetermined intervals, both random numbers stored in block key section 502 and data section 504 are transformed. The random number provided to block key section 502 is independently

transformed, whereas transformation of the random number provided to data section 504 is dependent on the transformation being performed in block key section 502. Mapping block 506 retrieves a subset each, of the newly transformed states of the two random numbers, and reduces them to generate one bit of the pseudo random bit sequence. Thus, in a desired number of clock cycles, a pseudo random bit sequence of a desired length is generated.

For the illustrated embodiment, by virtue of the employment of the "rekeying enable" signal, stream key section 506 may be left operating even during the block mode, as its outputs are effectively discarded by the "rekeying enable" signal (set in a "disabled" state).

Figure 6 illustrates the block key section of Fig. 5 in further detail, in accordance with one embodiment. As illustrated, block key section 502 includes registers 602a-602c, substitution boxes 604, and linear transformation unit 606. In block mode, registers 602a-602c are collectively initialized to a block cipher key, e.g. authentication key K_m or session key K_s . In stream mode, registers 602a-602c are collectively initialized to a random number, e.g. session key K_s or frame key K_i . Each round, substitution boxes 604 and linear transformation unit 606 modify the content of registers 602a-602c. More specifically, substitution boxes 604 receive the content of register 602a, modify it, and then store the substituted content into register 602c. Similarly, linear transformation unit 606 receives the content of registers 602b and 602c, linearly transforms them, and then correspondingly stores the linearly transformed content into registers 602a and 602b.

Substitution boxes 604 and linear transformation unit 606 may be implemented in a variety of ways in accordance with well known cryptographic principles. One specific implementation is given in more detail below after the description of Fig. 7.

Figure 7 illustrates the block data section of Fig. 5 in further detail, in accordance with one embodiment. For the illustrated embodiment, data section 504 is similarly constituted as block key section 502, except linear transformation unit 706 also takes into consideration the content of register 602b, when

transforming the contents of registers **702b-702c**. In block mode, registers **702a-702c** are collectively initialized with the target plain text, e.g. earlier described random number A_n or derived random number M_{i-1} . In stream mode, registers **702a-702c** are collectively initialized with a random number. Each round, substitution boxes **704** and linear transformation unit **706** modify the content of registers **702a-702c** as described earlier for block key section **502** except for the differences noted above.

Again, substitution boxes **604** and linear transformation unit **606** may be implemented in a variety of ways in accordance with well known cryptographic principles.

In one implementation for the above described embodiment, each register **602a**, **602b**, **602c**, **702a**, **702b**, **702c** is 28-bit wide. [Whenever registers **602a-602c** or **702a-702c** collectively initialized with a key value or random number less than 84 bits, the less than 84-bit number is initialized to the lower order bit positions with the higher order bit positions zero filled.] Additionally, each set of substitution boxes **604** or **704** are constituted with seven 4 input by 4 output substitution boxes. Each linear transformation unit **606** or **706** produces 56 output values by combining outputs from eight diffusion networks (each producing seven outputs). More specifically, the operation of substitution boxes **604/704** and linear transformation unit **606/706** are specified by the four tables to follow. For substitution boxes **604/704**, the l th input to box J is bit $l*7+J$ of register **602a/702a**, and output l of box J goes to bit $l*7+j$ of register **602c/702c**. [Bit 0 is the least significant bit.] For each diffusion network (linear transformation unit **606** as well as **706**), the inputs are generally labeled $I0-I6$ and the outputs are labeled $O0-O6$. The extra inputs for each diffusion network of the linear transformation unit **706** is labeled $K0-K6$.

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
SK	8	14	5	9	3	0	12	6	1	11	15	2	4	7	10	13
SK	1	6	4	15	8	3	11	5	10	0	9	12	7	13	14	2
SK	13	11	8	6	7	4	2	15	1	12	14	0	10	3	9	5
SK	0	14	11	7	12	3	2	13	15	4	8	1	9	10	5	6
SK	12	7	15	8	11	14	1	4	6	10	3	5	0	9	13	2

SK	1	12	7	2	8	3	4	14	11	5	0	15	13	6	10	9
SK	10	7	6	1	0	14	3	13	12	9	11	2	15	5	4	8
SB	12	9	3	0	11	5	13	6	2	4	14	7	8	15	1	10
SB	3	8	14	1	5	2	11	13	10	4	9	7	6	15	12	0
SB	7	4	1	10	11	13	14	3	12	15	6	0	2	8	9	5
SB	6	3	1	4	10	12	15	2	5	14	11	8	9	7	0	13
SB	3	6	15	12	4	1	9	2	5	8	10	7	11	13	0	14
SB	11	14	6	8	5	2	12	7	1	4	15	3	10	13	9	0
SB	1	11	7	4	2	5	12	9	13	6	8	15	14	0	3	10

Table I – Substitution performed by each of the seven constituting substitution boxes of substitution boxes 604/704.

	Diffusion Network Logic Function
O_0	$K_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$
O_1	$K_1 \oplus I_0 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$
O_2	$K_2 \oplus I_0 \oplus I_1 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$
O_3	$K_3 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_4 \oplus I_5 \oplus I_6$
O_4	$K_4 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_5 \oplus I_6$
O_5	$K_5 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_6$
O_6	$K_6 \oplus I_0 \oplus I_1 \oplus I_2 \oplus I_3 \oplus I_4 \oplus I_5 \oplus I_6$

	K1	K2	K3	K4	K5	K6	K7	K8
I_0	Kz0	Ky0	Ky4	Ky8	Ky12	Ky16	Ky20	Ky24

I ₁	Kz1	Ky1	Ky5	Ky9	Ky13	Ky17	Ky21	Ky25
I ₂	Kz2	Ky2	Ky6	Ky10	Ky14	Ky18	Ky22	Ky26
I ₃	Kz3	Ky3	Ky7	Ky11	Ky15	Ky19	Ky23	Ky27
I ₄	Kz4	Kz7	Kz10	Kz13	Kz16	Kz19	Kz22	Kz25
I ₅	Kz5	Kz8	Kz11	Kz14	Kz17	Kz20	Kz23	Kz26
I ₆	Kz6	Kz9	Kz12	Kz15	Kz18	Kz21	Kz24	Kz27
O ₀	Kx0	Ky0	Ky1	Ky2	Ky3	Kx7	Kx8	Kx9
O ₁	Kx1	Ky4	Ky5	Ky6	Ky7	Kx10	Kx11	Kx12
O ₂	Kx2	Ky8	Ky9	Ky10	Ky11	Kx13	Kx14	Kx15
O ₃	Kx3	Ky12	Ky13	Ky14	Ky15	Kx16	Kx17	Kx18
O ₄	Kx4	Ky16	Ky17	Ky18	Ky19	Kx19	Kx20	Kx21
O ₅	Kx5	Ky20	Ky21	Ky22	Ky23	Kx22	Kx23	Kx24
O ₆	Kx6	Ky24	Ky25	Ky26	Ky27	Kx25	Kx26	Kx27

Tables II & III – Diffusion networks for linear transformation unit 606/706

(continued in Table IV).

	B1	B2	B3	B4	B5	B6	B7	B8
I ₀	Bz0	By0	By4	By8	By12	By16	By20	By24
I ₁	Bz1	By1	By5	By9	By13	By17	By21	By25
I ₂	Bz2	By2	By6	By10	By14	By18	By22	By26
I ₃	Bz3	By3	By7	By11	By15	By19	By23	By27
I ₄	Bz4	Bz7	Bz10	Bz13	Bz16	Bz19	Bz22	Bz25
I ₅	Bz5	Bz8	Bz11	Bz14	Bz17	Bz20	Bz23	Bz26
I ₆	Bz6	Bz9	Bz12	Bz15	Bz18	Bz21	Bz24	Bz27
K ₀	Ky0	–	–	–	–	Ky7	Ky14	Ky21
K ₁	Ky1	–	–	–	–	Ky8	Ky15	Ky22
K ₂	Ky2	–	–	–	–	Ky9	Ky16	Ky23

K_3	Ky3	–	–	–	–	Ky10	Ky17	Ky24
K_4	Ky4	–	–	–	–	Ky11	Ky18	Ky25
K_5	Ky5	–	–	–	–	Ky12	Ky19	Ky26
K_6	Ky6	–	–	–	–	Ky13	Ky20	Ky27
O_0	Bx0	By0	By1	By2	By3	Bx7	Bx8	Bx9
O_1	Bx1	By4	By5	By6	By7	Bx10	Bx11	Bx12
O_2	Bx2	By8	By9	By10	By11	Bx13	Bx14	Bx15
O_3	Bx3	By12	By13	By14	By15	Bx16	Bx17	Bx18
O_4	Bx4	By16	By17	By18	By19	Bx19	Bx20	Bx21
O_5	Bx5	By20	By21	By22	By23	Bx22	Bx23	Bx24
O_6	Bx6	By24	By25	By26	By27	Bx25	Bx26	Bx27

Table IV – Diffusion networks for linear transformation unit 606/706

(continued from Tables II & III).

Referring now back to **Fig. 5**, recall that a ciphered text may be deciphered by generating the intermediate “keys” and applying them backward. Alternatively, for an embodiment where either the inverse of substitution boxes 604/704 and linear transformation units 606/706 are included or they may be dynamically reconfigured to operate in an inverse manner, the ciphered text may be deciphered as follows. First, the cipher key used to cipher the plain text is loaded into block key section 502, and block key section 502 is advanced by R-1 rounds, i.e. one round short of the number of rounds (R) applied to cipher the plain text. After the initial R-1 rounds, the ciphered text is loaded into data section 504, and both sections, block key section 502 and data section 504, are operated “backward”, i.e. with substitution boxes 604/704 and linear transformation units 606/706 applying the inverse substitutions and linear transformations respectively.

Figures 8a-8c illustrate the stream key section of **Fig. 5** in further detail, in accordance with one embodiment. As illustrated in **Fig. 8a**, stream key section 506 includes a number of linear feedback shift registers (LFSRs) 802 and

combiner function **804**, coupled to each other as shown. LFSRs **802** are collectively initialized with a stream cipher key, e.g. earlier described frame key K_i . During operation, the stream cipher key is successively shifted through LFSRs **802**. Selective outputs are taken from LFSRs **802**, and combiner function **804** is used to combine the selective outputs. In stream mode (under which, rekeying is enabled), the combined result is used to dynamically modify a then current state of a block cipher key in block key section **502**.

For the illustrated embodiment, four LFSRs of different lengths are employed. Three sets of outputs are taken from the four LFSRs. The polynomials represented by the LFSR and the bit positions of the three sets of LFSR outputs are given by the table to follows:

LFSR	Polynomial	Combining Function		
		Taps		
		0	1	2
3	$X^{17} + X^{15} + X^{11} + X^5 + 1$	6	12	17
2	$X^{16} + X^{15} + X^{12} + X^8 + X^7 + X^5 + 1$	6	10	16
1	$X^{14} + X^{11} + X^{10} + X^7 + X^6 + X^4 + 1$	5	9	14
0	$X^{13} + X^{11} + X^9 + X^5 + 1$	4	8	13

Table V – Polynomials of the LFSR and tap positions.

The combined result is generated from the third set of LFSR outputs, using the first and second set of LFSR outputs as data and control inputs respectively to combiner function **802**. The third set of LFSR outputs are combined into a single bit. In stream mode (under which, rekeying is enabled), the combined single bit is then used to dynamically modify a predetermined bit of a then current state of a block cipher key in block key section **502**.

Fig. 8b illustrates combiner function **804** in further detail, in accordance with one embodiment. As illustrated, combiner function **804** includes shuffle network **806** and XOR **808a-808b**, serially coupled to each other and LFSRs **802** as shown. For the illustrated embodiment, shuffle network **806** includes four binary shuffle units **810a-810d** serially coupled to each other, with first and last binary shuffle units **810a** and **810d** coupled to XOR **808a** and **808b** respectively. XOR **808a** takes the first group of LFSR outputs and combined them as a single bit input for shuffle network **806**. Binary shuffle units **810a-810d** serially propagate and shuffle the output of XOR **808a**. The second group of LFSR outputs are used to control the shuffling at corresponding ones of binary shuffle units **810a-810d**. XOR **808b** combines the third set of LFSR outputs with the output of last binary shuffle unit **810d**.

Fig. 8c illustrates one binary shuffle unit **810*** (where * is one of a-d) in further detail, in accordance with one embodiment. Each binary shuffle unit **810*** includes two flip-flops **812a** and **812b**, and a number of selectors **814a-814c**, coupled to each other as shown. Flip-flops **812a** and **812b** are used to store two state values (A, B). Each selector **814a**, **814b** or **814c** receives a corresponding one of the second group of LFSR outputs as its control signal. Selector **814a-814b** also each receives the output of XOR **808a** or an immediately preceding binary shuffle unit **810*** as input. Selector **814a-814b** are coupled to flip-flops **812a-812b** to output one of the two stored state values and to shuffle as well as modify the stored values in accordance with the state of the select signal. More specifically, for the illustrated embodiment, if the stored state values are (A, B), and the input and select values are (D, S), binary shuffle unit **810*** outputs A, and stores (B, D) if the value of S is "0". Binary shuffle unit **810*** outputs B, and stores (D, A) if the value of S is "1".

Referring now to back to **Figure 5**, as illustrated and described earlier, mapping function **508** generates the pseudo random bit sequence based on the contents of selected registers of block key section **502** and data section **504**. In one embodiment, where block key section **502** and data section **504** are implemented in accordance with the respective embodiments illustrated in **Fig. 6-**

7, mapping function 508 generates the pseudo random bit sequence at 24-bit per clock based on the contents of registers (Ky and Kz) 602b-602c and (By and Bz) 702b-702c. More specifically, each of the 24 bits is generated by performing the XOR operation on nine terms in accordance with the following formula:

$$(B0 \cdot K0) \oplus (B1 \cdot K1) \oplus (B2 \cdot K2) \oplus (B3 \cdot K3) \oplus (B4 \cdot K4) \oplus (B5 \cdot K5) \oplus (B6 \cdot K6) \oplus B7 \oplus K7$$

Where " \oplus " represents a logical XOR function, " \cdot " represents a logical AND function, and the input values B and K for the 24 output bits are

Input	B0	B1	B2	B3	B4	B5	B6	B7	K0	K1	K2	K3	K4	K5	K6	K7
Origin	Bz	Bz	Bz	Bz	Bz	Bz	Bz	By	Kz	Kz	Kz	Kz	Kz	Kz	Kz	Ky
Output bit																
	14	23	7	27	3	18	8	20	12	24	0	9	16	7	20	13
	20	26	6	15	8	19	0	10	26	18	1	11	6	20	12	19
	7	20	2	10	19	14	26	17	1	22	8	13	7	16	25	3
	22	12	6	17	3	10	27	4	24	2	9	5	14	18	21	15
	22	24	14	18	7	1	9	21	19	24	20	8	13	6	3	5
	12	1	16	5	10	24	20	14	27	2	8	16	15	22	4	21
	5	3	27	8	17	15	21	12	14	23	16	10	27	1	7	17
	9	20	1	16	5	25	12	6	9	13	22	17	1	24	5	11
	23	25	11	13	17	1	6	22	25	21	18	15	6	11	1	10
	4	0	22	17	25	10	15	18	0	20	26	19	4	15	9	27
1	23	25	9	2	13	16	4	8	2	11	27	19	14	22	4	7
1	3	6	20	12	25	19	10	27	24	3	14	6	23	17	10	1
1	26	1	18	21	14	4	10	0	17	7	26	0	23	11	14	8
1	2	11	4	21	15	24	18	9	5	16	12	2	26	23	11	6
1	22	24	3	19	11	4	13	5	22	0	18	8	25	5	15	2
1	12	0	27	11	22	5	16	1	10	3	15	19	21	27	6	18
1	24	20	2	7	15	18	8	3	12	20	5	19	1	27	8	23
1	12	16	8	24	7	2	21	23	17	2	11	14	7	25	22	16
1	19	3	22	9	13	6	25	7	4	10	2	17	21	24	13	22
1	11	17	13	26	4	21	2	16	3	4	13	26	18	23	9	25
2	17	23	26	14	5	11	0	15	26	3	9	19	21	12	6	0
2	9	14	23	16	27	0	6	24	18	21	3	27	4	10	15	26

2	7	21	8	13	1	26	19	25	25	0	12	10	7	17	23	9
2	27	15	23	5	0	9	18	11	8	0	25	20	16	5	13	12

Accordingly, a novel method and apparatus for ciphering and deciphering video content to protect the video content from unauthorized copying during transmission has been described.

Epilogue

From the foregoing description, those skilled in the art will recognize that many other variations of the present invention are possible. In particular, while the present invention has been described as being implemented in interfaces **108a** and **108b**, some of the logic may be distributed in other components of video source and sink devices **102** and **104**. Additionally, non-LFSR based stream key section, more or less block key registers, larger or smaller block key registers, more or less substitution units, including alternative substitution patterns, as well as different linear transformation units may be employed. Thus, the present invention is not limited by the details described, instead, the present invention can be practiced with modifications and alterations within the spirit and scope of the appended claims.

CLAIMS

What is claimed is:

1. In a video source device, a method comprising:
providing a basis value to a symmetric ciphering/deciphering process to a video sink device to receive a video content from the video source device;
ciphering the video content for transmission to the video sink device,
including generation of a first cipher key through functional transformation of the basis value; and
verifying periodically that the transmitted video content is indeed being symmetrically deciphered by the video sink device.
2. The method of claim 1, wherein said ciphering process comprises subsequent generation of a plurality of cipher keys using the initial cipher key.
3. The method of claim 1, wherein said functional transformation of the basis value is being performed employing an authentication key generated as an integral part of an authentication process authenticating the video sink device.
4. The method of claim 3, wherein said authentication key is generated using identifiers of the video source and sink devices, and said authentication process include exchanging the identifiers with the video sink device.
5. The method of claim 1, wherein said periodic verification comprises receiving selected ones of a plurality of verification values from the video sink device at predetermined points in time during said ciphering/deciphering process, and comparing the received verification values to corresponding ones of a plurality of verification reference values.
6. The method of claim 5, wherein said ciphering process comprises generating a pseudo random bit sequence for ciphering the video content, said verification reference values are symmetrically derived from the pseudo random bit sequence.
7. In a video source device, a method comprising:
authenticating a video sink device to receive a video content from the video source device, including generation of an authentication key;
ciphering the video content for transmission to the video sink device,
including generation of an initial cipher key using at least the authentication key;
and
verifying periodically the transmitted video content is being symmetrically deciphered by the video sink device.
8. The method of claim 7, wherein said ciphering process further comprises subsequent generation of a plurality of cipher keys using the initial cipher key.

9. The method of claim 7, wherein said authentication key is generated using identifiers of the video source and sink devices, and said authentication process include exchanging the identifiers with the video sink device.
10. The method of claim 7, wherein said periodic verification comprises receiving selected ones of a plurality of verification values from the video sink device at predetermined points in time during the ciphering/deciphering process, and comparing the received verification values to corresponding ones of a plurality of verification reference values.
11. The method of claim 10, wherein said ciphering process comprises generating a pseudo random bit sequence for ciphering the video content, said verification reference values are symmetrically derived from the pseudo random bit sequence.
12. An apparatus comprising:
a first control unit to provide a basis value to a symmetric ciphering/deciphering process to a video sink device to receive a video content from the video source device;
a cipher unit coupled to the first control unit to cipher the video content for transmission to the video sink device, wherein the ciphering unit generates a first cipher key through functional transformation of the basis value; and
a second control unit coupled to the cipher unit to periodically verify that the transmitted video content is indeed being symmetrically deciphered by the video sink device.
13. The apparatus of claim 12, wherein said cipher unit comprises a block cipher to perform said functionally transformation of the basis value to generate the initial cipher key, and to successively generate a plurality of additional cipher keys using the initial cipher key.
14. The apparatus of claim 13, wherein said first control unit also authenticates the video sink device and generates an authentication key as an integral part of the authentication process, and the block cipher functionally transforms the basis value employing the authentication key.
15. The apparatus of claim 12, wherein said second control unit receives selected ones of a plurality of verification values from the video sink device at predetermined points in time during said ciphering/deciphering process, and compares the received verification values to corresponding ones of a plurality of verification reference values.
16. The apparatus of claim 15, wherein said cipher unit comprises a stream cipher to generate a pseudo random bit sequence for ciphering the video content, and said second control unit derives said verification reference values from the pseudo random bit sequence.
17. An apparatus comprising:

a first control unit to authenticate a video sink device to receive a video content from the video source device, including generation of an authentication key;

a cipher unit coupled to the first control unit to cipher the video content for transmission to the video sink device, including generation of an initial cipher key using at least the authentication key; and

a second control unit coupled to the cipher unit to verify periodically the transmitted video content is being symmetrically deciphered by the video sink device.

18. The apparatus of claim 17, wherein said cipher unit comprises a block cipher to generate said first cipher key using at least said authentication key, and to successively generate a plurality of additional cipher keys using the initial cipher key.

19. The apparatus of claim 17, wherein said first control unit also exchanges device identifiers with the video sink device, and generates the authentication key using the device identifiers.

20. The apparatus of claim 17, wherein said second control unit receives selected ones of a plurality of verification values from the video sink device at predetermined points in time during the ciphering/deciphering process, and compares the received verification values to corresponding ones of a plurality of verification reference values.

21. The apparatus of claim 20, wherein said cipher unit comprises a stream cipher to generate a pseudo random bit sequence for ciphering/deciphering the video content, said verification reference values are derived from the pseudo random bit sequence.

22. In a video sink device, a method comprising:
receiving a basis value to a symmetric ciphering/deciphering process from a video source device to provide a video content in ciphered form to the video sink device;

deciphering received video content from the video source device, including generation of a first decipher key through functional transformation of the basis value; and

providing periodically verification values to the video source device to facilitate confirmation that the ciphered video content is indeed being symmetrically deciphered.

23. The method of claim 22, wherein said deciphering process comprises subsequent generation of a plurality of decipher keys using the initial decipher key.

24. The method of claim 22, wherein said functional transformation of the basis value is being performed employing an authentication key generated as an integral part of an authentication process authenticating the video source device.

25. The method of claim 24, wherein said authentication key is generated using identifiers of the video source and sink devices, and said authentication process include exchanging the identifiers with the video source device.
26. The method of claim 22, wherein said deciphering processes comprises generating a pseudo random bit sequence for deciphering the video content, said periodic provision of selected ones of verification values include symmetrically deriving the verification values from the pseudo random bit sequence.
27. In a video sink device, a method comprising:
 authenticating a video source device to provide a video content in ciphered form to the video sink device, including generation of an authentication key;
 deciphering the ciphered video content received from the video source device, including generation of an initial decipher key using at least the authentication key; and
 providing periodically verification values to the video source device to facilitate confirmation that the ciphered video content is indeed being symmetrically deciphered.
28. The method of claim 27, wherein said deciphering process further comprises subsequent generation of a plurality of decipher keys using the initial decipher key.
29. The method of claim 27, wherein said authentication key is generated using identifiers of the video source and sink devices, and said authentication process include exchanging the identifiers with the video source device.
30. The method of claim 27, wherein said deciphering process comprises generating a pseudo random bit sequence for deciphering the video content, and said verification values are symmetrically derived from the pseudo random bit sequence.
31. An apparatus comprising:
 a first control unit to receive a basis value to a symmetric ciphering/deciphering process from a video source device to provide a video content in ciphered form to the apparatus;
 a decipher unit coupled to the first control unit to decipher the ciphered video content, wherein the decipher unit generates a first decipher key through functional transformation of the basis value; and
 a second control unit coupled to the decipher unit to periodically provide verification values to the video source device to facilitate confirmation that indeed the ciphered video content is being symmetrically deciphered.
32. The apparatus of claim 31, wherein said decipher unit comprises a block cipher to perform said functionally transformation of the basis value to generate

the initial decipher key, and to successively generate a plurality of additional decipher keys using the initial decipher key.

33. The apparatus of claim 32, wherein said first control unit also authenticates the video source device and generates an authentication key as an integral part of the authentication process, and the block cipher functionally transforms the basis value employing the authentication key.

34. The apparatus of claim 32, wherein said decipher unit comprises a stream cipher to generate a pseudo random bit sequence for deciphering the video content, and said second control unit derives said verification values from the pseudo random bit sequence.

35. An apparatus comprising:

- a first control unit to authenticate a video source device to provide a video content in ciphered form to the apparatus, including generation of an authentication key;

- a decipher unit coupled to the first control unit to decipher the ciphered video content, including generation of an initial decipher key using at least the authentication key; and

- a second control unit coupled to the decipher unit to provide periodically verification values to the video source device to facilitate confirmation that the ciphered video content is indeed being symmetrically deciphered.

36. The apparatus of claim 35, wherein said decipher unit comprises a block cipher to generate said first decipher key using at least said authentication key, and to successively generate a plurality of additional decipher keys using the initial decipher key.

37. The apparatus of claim 35, wherein said first control unit also exchanges device identifiers with the video source device, and generates the authentication key using the device identifiers.

38. The apparatus of claim 35, wherein said decipher unit comprises a stream cipher to generate a pseudo random bit sequence for deciphering the video content, said verification values are derived from the pseudo random bit sequence.

1/5

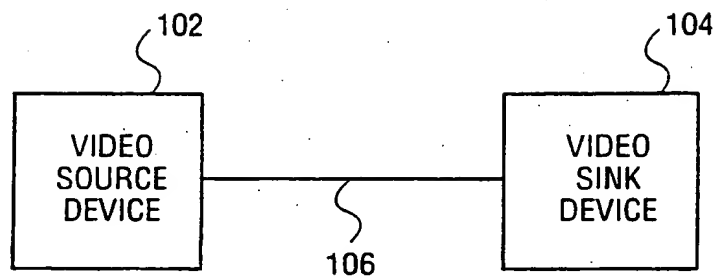


FIG. 1

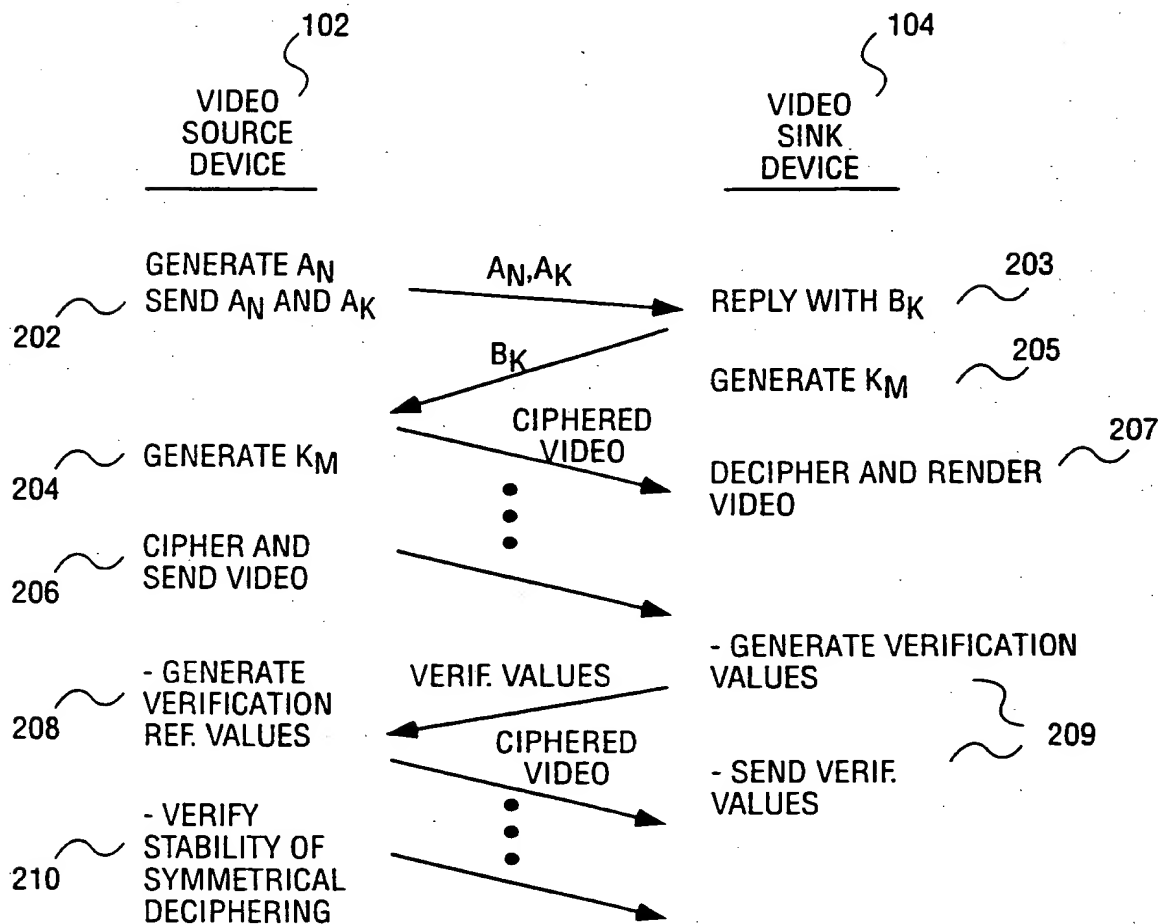
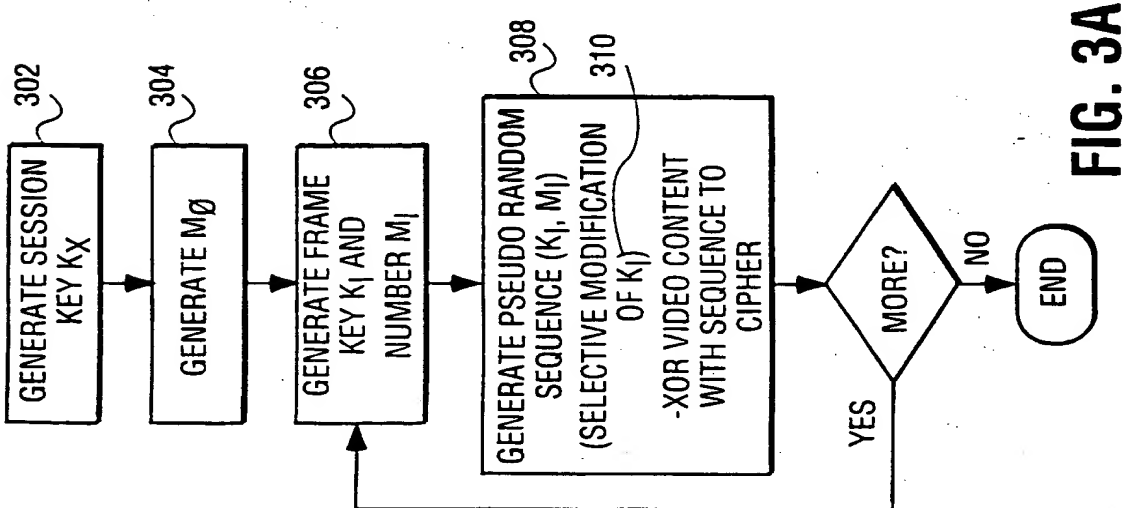
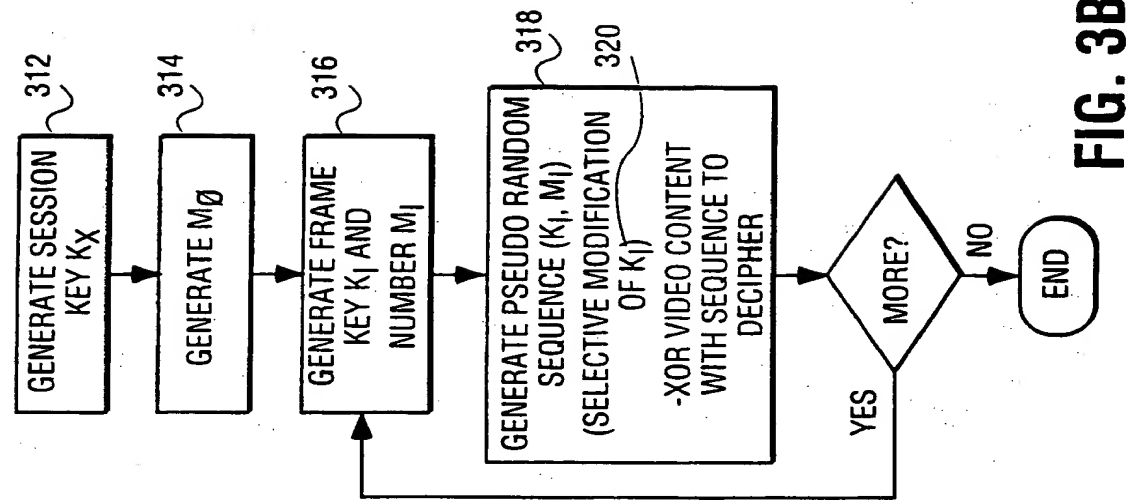


FIG. 2

THIS PAGE BLANK (USPTO)



THIS PAGE BLANK (USPTO)

3/5

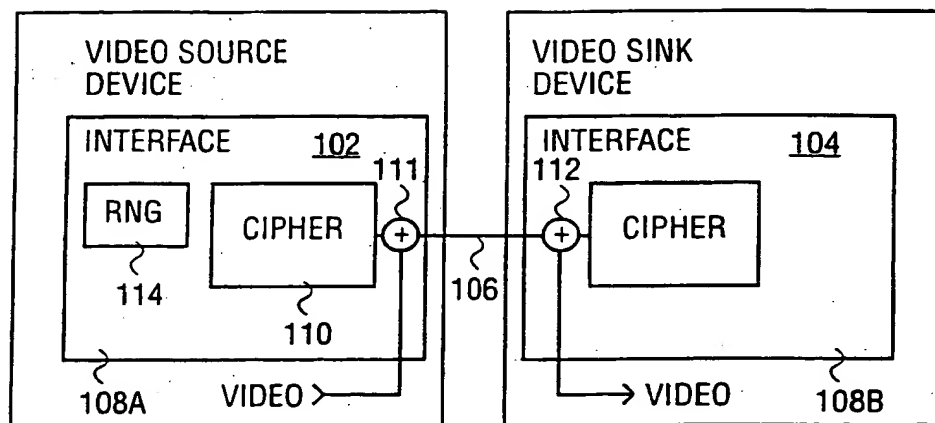


FIG. 4

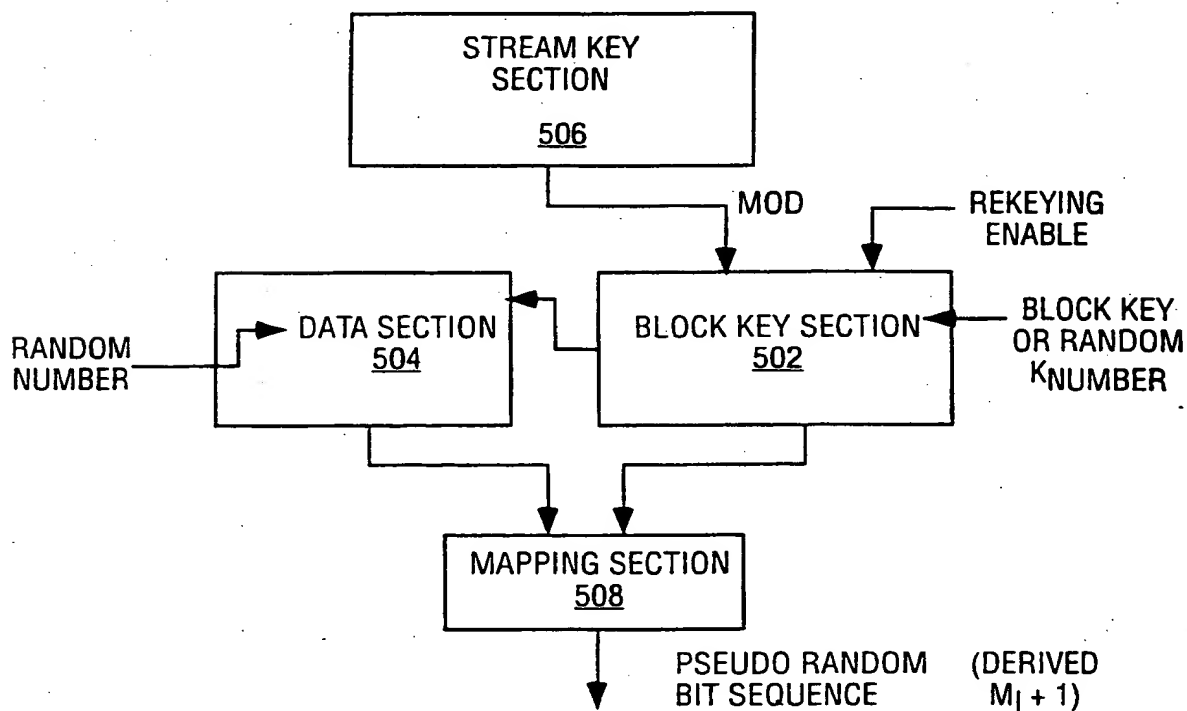
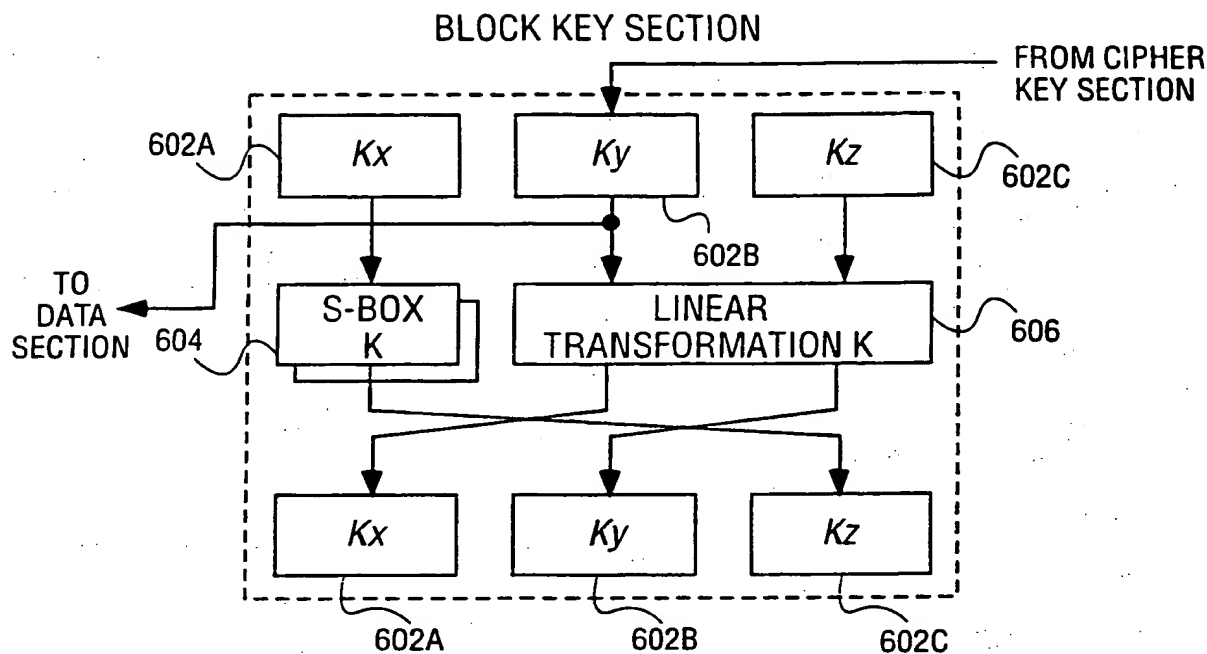
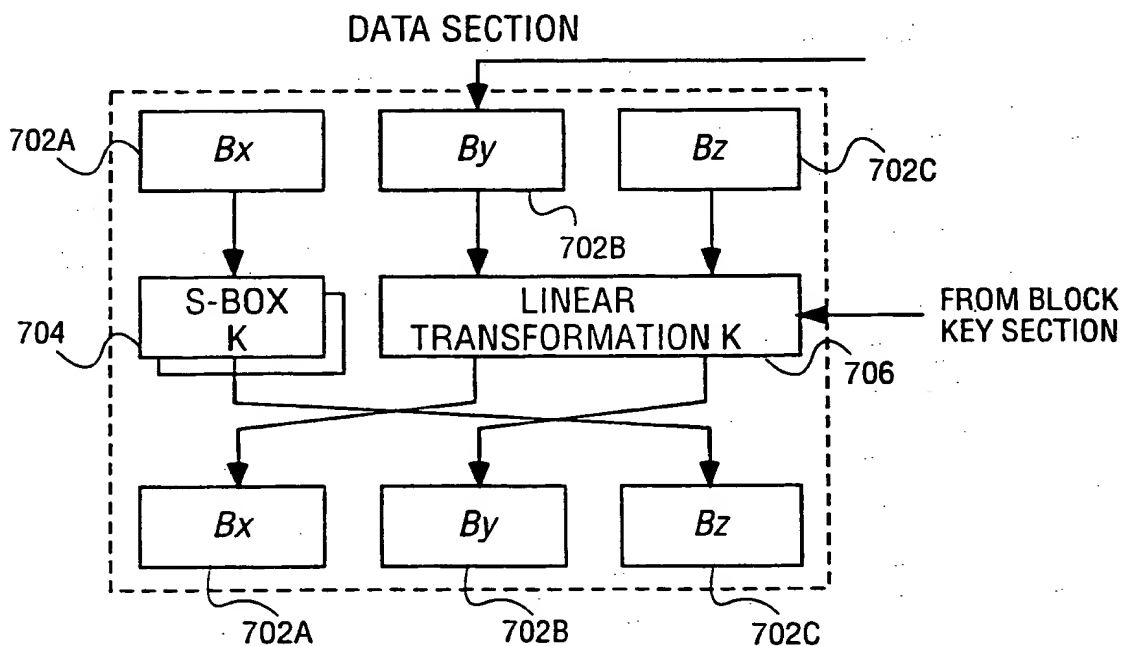


FIG. 5

THIS PAGE BLANK (USPTO)

4/5

**FIG. 6****FIG. 7**

THIS PAGE BLANK (USPTO)

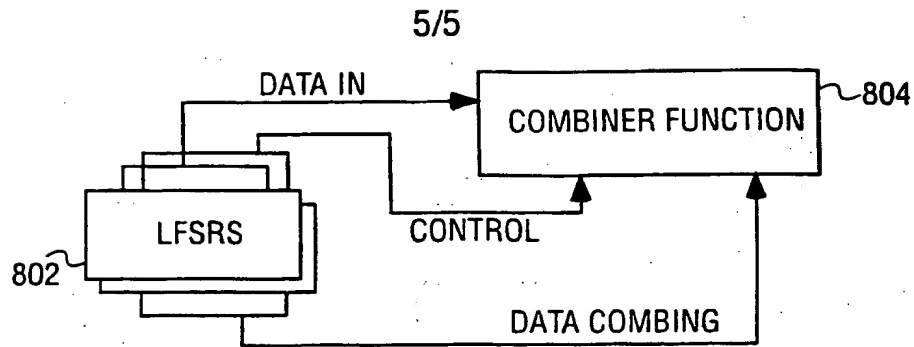


FIG. 8A

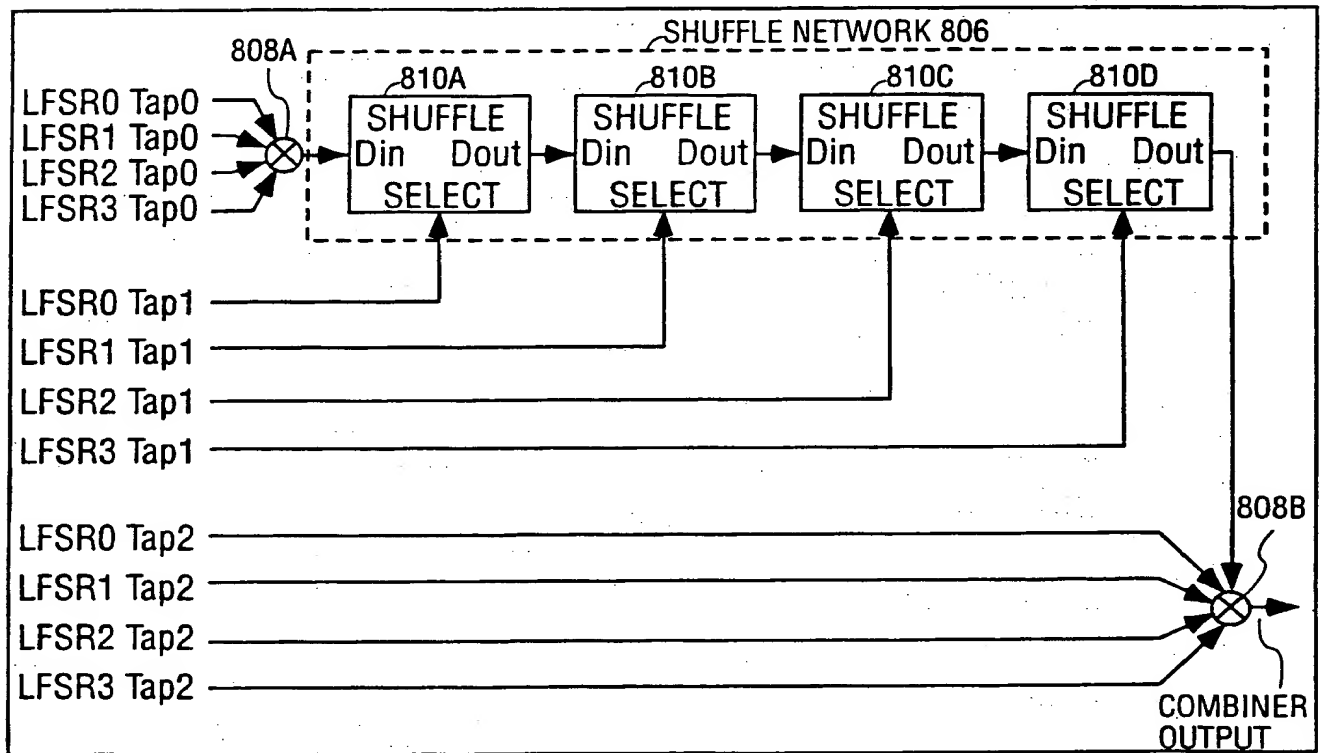


FIG. 8B

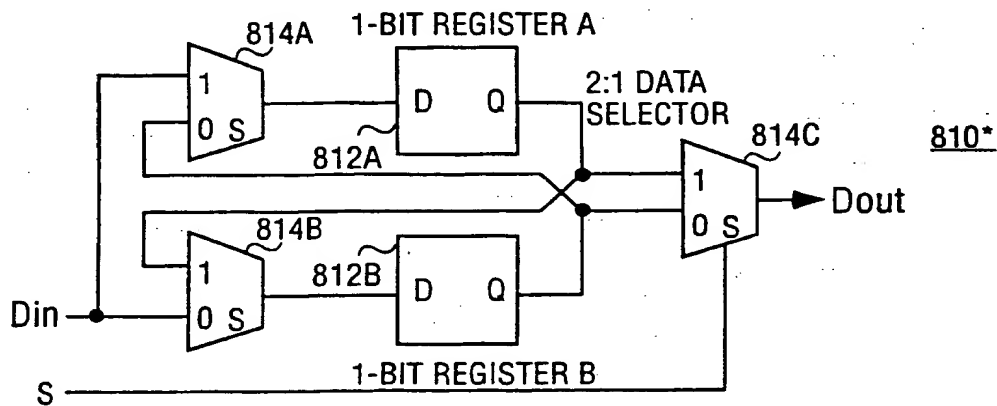


FIG. 8C

THIS PAGE BLANK (USPTO)

INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 00/22834

A. CLASSIFICATION OF SUBJECT MATTER
IPC 7 H04N7/167

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 7 H04N

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

EP0-Internal

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	WO 99 19822 A (MICROSOFT CORP) 22 April 1999 (1999-04-22) page 5, line 11 -page 11, line 5 page 12, line 23 -page 17, line 19 figures 1-5	1-38
A	US 5 852 472 A (GUTMANN MICHAEL J ET AL) 22 December 1998 (1998-12-22) column 1, line 14 -column 29 column 3, line 4 -column 6, line 42 figures 3,4	1-38

☐ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

*** Special categories of cited documents :**

A document defining the general state of the art which is not considered to be of particular relevance

E earlier document but published on or after the international filing date

L document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

O document referring to an oral disclosure, use, exhibition or other means

P document published prior to the international filing date but later than the priority date claimed

T later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

X document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

Y document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

8 document member of the same patent family

Date of the actual completion of the international search

28 November 2000

Date of mailing of the international search report

04/12/2000

Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2
NL - 2280 HV Rijswijk
Tel. (+31-70) 340-2040. Tx. 31 651 epo nl.
Fax: (+31-70) 340-3016

Authorized officer

Van der Zaal, R

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No

PCT/US 00/22834

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
WO 9919822 A	22-04-1999	EP 1031206 A	30-08-2000
US 5852472 A	22-12-1998	NONE	